

Programming Languages [Fall 2019] Test III

NAME: _____

Instructions:

- 1) This test is 10 pages in length.
- 2) You have 2 hours to complete and turn in this test.
- 3) This test is closed books, notes, laptops, phones, smartwatches, friends, neighbors, etc.
- 4) Use the backs of pages in this test packet for scratch work.

- 5) Write and sign the following: "I pledge my Honor that I have not cheated, and will not cheat, on this test."

Signed: _____

1. [4 points]

Define a monomorphic function in ML. For full credit, define the function using as few characters as possible.

2. [4 points]

How does ML avoid having a NULL value? [2-3 sentences]

3. [4 points]

For each of the following ML expressions, write the expression's type, or, if the expression is ill typed, briefly explain why.

a) `fun f(x) = if f(x) then f else (fn x=>x)`

b) `fun f(x) = if true then (fn x=>if x<1 then nil else nil) else (fn x=>[true])`

[Undergraduates skip the following problem. Problem 4 is for graduate students.]

4. [5 points]

Explain the ML value restriction. Then define functionally equivalent ML expressions e and e' such that e violates the value restriction but e' does not.

5. [12 points]

An ML API has the limited documentation shown below. Although limited, the documentation provides enough information to understand and implement the functions. Implement each of these two functions *twice*, first according to the constraints of Assignment 1 (i.e., with recursion) and second according to the constraints of Assignment 2 (i.e., with map/fold).

a) `val find : ('a -> bool)->'a list->'a option (*finds leftmost match*)`

b) `val partition : ('a -> bool) -> 'a list -> ('a list * 'a list)`
`(*the returned pair has the true-returning elements in the first list*)`
`(*elements in the returned lists are ordered as in the argument list*)`

6. [8 points]

Prove the standard type-safety corollary, assuming that Progress and Preservation have been already proved.

7. [10 points]

Building on our encodings of natural numbers as Church numerals, encode a multiplication operator into λ_{UT} . Then trace execution of 2×3 with the normal-order evaluation strategy. As we did in class, underline redexes and use abbreviations (e.g., by writing 2 instead of the full lambda expression encoding 2).

8. [10 points]

Using evaluation contexts, define dynamic semantics for the simply typed lambda calculus having base type unit. Define two versions of the dynamic semantics: first with call-by-name evaluation and second with full-beta evaluation.

9. [13 points]

a) Define alpha-equivalence for the simply typed lambda calculus having base type unit. Assume that definitions of free variables and capture-avoiding substitution already exist. As in class, do not define explicit rules for symmetry or transitivity.

b) Prove that your definition of alpha-equivalence is in fact symmetric.
Theorem. For all expressions e and e' : if $e \equiv e'$ then $e' \equiv e$.

10. [12 points]

Let S be the call-by-value simply typed lambda calculus with base type unit. State the Progress and Preservation theorems for S , and for each theorem, prove one case—the one case in which the beta-step rule is analyzed or used. Make all the assumptions we made in class, including that all the standard helper lemmas have already been proved.

11. [10 points]

Let D be diML with constructs added—as defined in class—for *binary* product and sum, unit, and recursive types. (Recall that D uses *let* expressions to eliminate product types and *case* expressions to eliminate sum types.) Implement in D a function `rev` that takes an arbitrary list L of integers and returns L reversed. In your response, please feel free to define and use abbreviations, but avoid non- D syntax (e.g., refrain from using syntactic sugar). Hint: Mimic our implementation of list-reversal in ML using difference lists.

12. [13 points]

Language U is the call-by-value simply typed lambda calculus having unit as the base type and having return expressions, as discussed in class. Define the first-order abstract syntax and complete static and dynamic semantics for U , as would be appropriate for a proof of type safety.